The general algebraic framework for Mathematical Fuzzy Logic

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Three stages of development of an area of logic

Chagrov (*K voprosu ob obratnoi matematike modal'noi logiki,* Online Journal Logical Studies, 2001)

distinguishes three stages in the development of a field in logic:

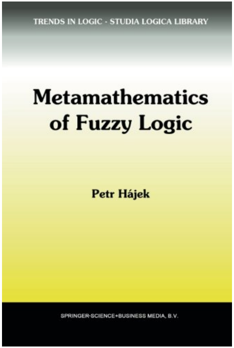
1. Emerging of the area
2. Development of particular logics and introduction of new ones
3. Universal methods

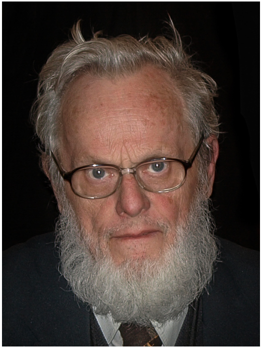
Three stages of development of MFL

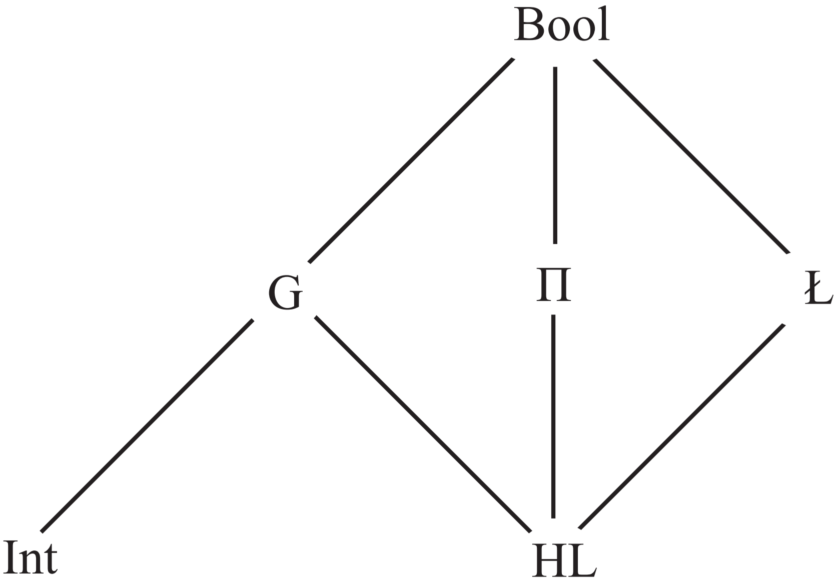
First stage: Emerging of the area (since 1965)

- 1965: Zadeh's fuzzy sets, 1968: 'fuzzy logic' (Goguen)
- 1970s: systems of fuzzy 'logic' lacking a good metatheory
- 1970s–1980s: first 'real' logics (Pavelka, Takeuti–Titani, . . .),
  discussion of many-valued logics in the fuzzy context

Culminated in Hájek's monograph (1998): G, Ł, Π, HL






Bool
G
Π
Ł
Int
HL

Three stages of development of MFL

Second stage: development of particular logics and introduction of many new ones (since the 1990s)

- New logics: MTL, SHL, UL, $\Pi_\sim$, ŁΠ, ...
- Algebraic semantics, proof theory, complexity

  Kripke-style and game-theoretic semantics, ...
- First-order, higher-order, and modal fuzzy logics

  Systematic treatment of particular fuzzy logics

Basic fuzzy logic?

Hájek called the logic HL the Basic fuzzy Logic BL

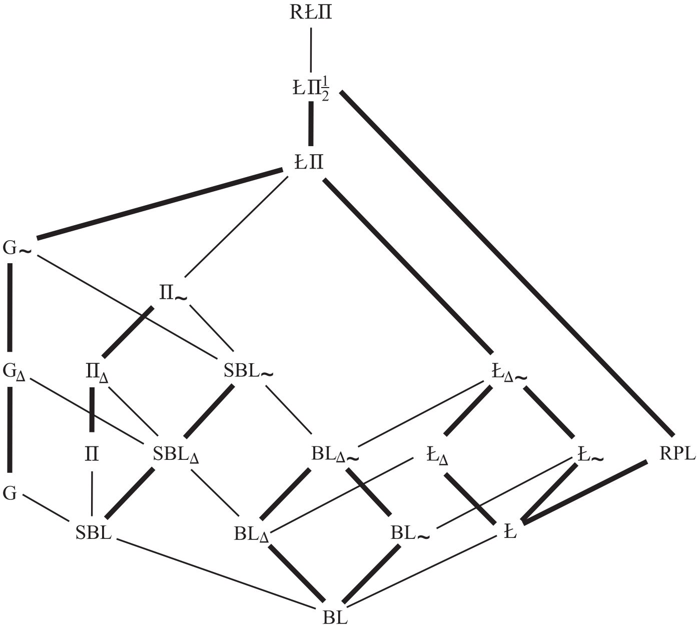
HL was *basic* in the following two senses:

1. *it could not be made weaker without losing essential properties*
2. *it provided a base for the study of all fuzzy logics.*

Because:

- HL is complete w.r.t. the semantics given by *all* continuous t-norms
- All fuzzy logics known by then were expansions of HL. The methods to introduce, algebraize, and study HL could be modified for all expansions of HL.

fuzzy logics = expansions of HL


RŁΠ
ŁΠ½
ŁΠ
G∼
Π∼
G∆
Π∆
SBL∼
Ł∆∼
Π
SBL∆
BL∆∼
Ł∆
Ł∼
RPL
G
SBL
BL∆
BL∼
Ł
BL

“Removing legs from the flea”

In the 3rd EUSFLAT (Zittau, Germany, September 2003) Petr Hájek started his lecture *Fleas and fuzzy logic: a survey* with a joke.

A group of scientists decide to investigate the ability of a flea can jump in relationship to how many legs it has.

They put the flea on a desk and said ’jump!’ The flea jumped and they noted:
“the flea with 6 legs can jump.”

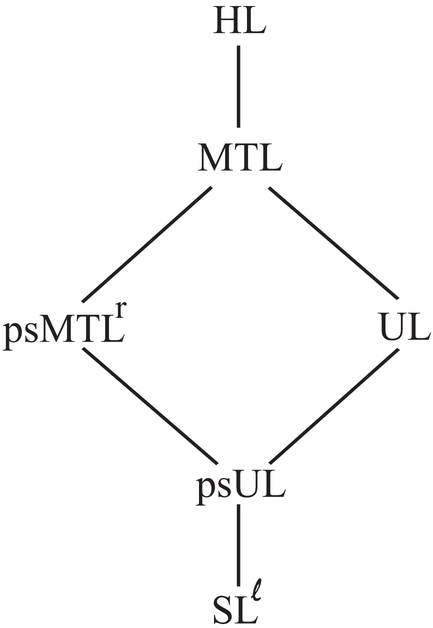
They remove a leg, repeated the command, the flea jumped and they noted:
“the flea with 5 legs can jump.”

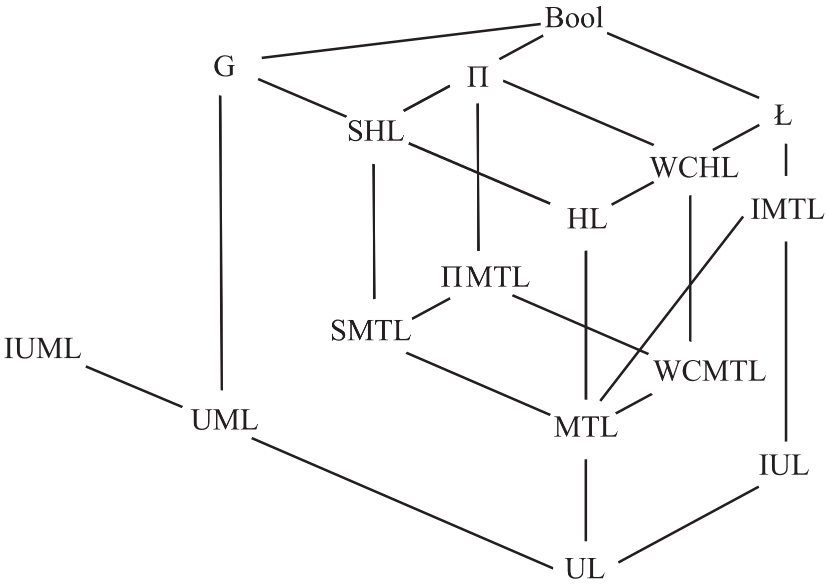
⋮

Finally, they removed the last legs repeated the command but
the flea didn’t move.

So they concluded:

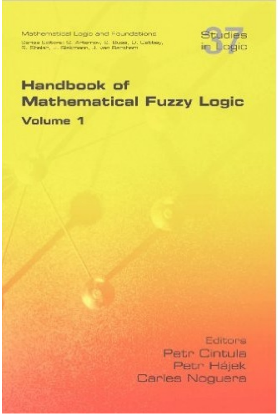
“Upon removing all its legs the flea loses sense of hearing.”

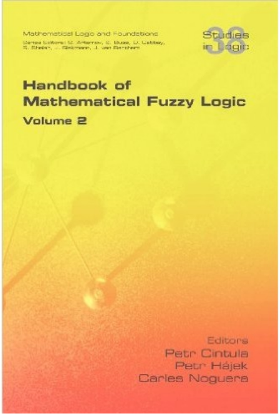

HL
MTL
psMTL^{r}
UL
psUL
SL$^{\ell}$

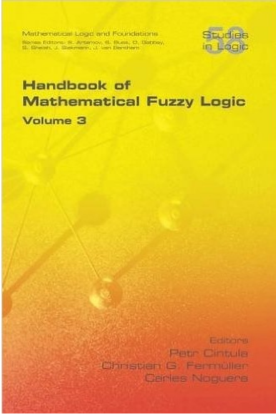

Bool
G
Π
Ł
SHL
WCHL
HL
IMTL
ΠMTL
SMTL
IUML
WCMTL
UML
MTL
IUL
UL

Three stages of development of MFL

The second stage is still ongoing; the state of the art is summarized in:







P. Cintula, C. Fermüller, P. Hájek, C. Noguera (editors). Vol. 37, 38, and 58 of *Studies in Logic: Math. Logic and Foundations*. College Publications, 2011, 2015.

Three stages of development of MFL

Third stage: universal methods (since $\sim$2006)

- There was a great deal of repetition in papers: slightly different logics studied by repeating the same definitions and essentially obtaining the same results by means of analogous proofs
- MFL needed general methods to prove metamathematical properties
- Classification of existing fuzzy logics
- Systematic treatment of classes of fuzzy logics
- Determining the position of fuzzy logics in the logical landscape

(Abstract) algebraic logic

- Algebraic logic: study of particular logical systems by giving them a semantics based on some algebraic structures
- Abstract algebraic logic (AAL): aims at understanding the various ways in which an arbitrary logical system can be endowed with an algebraic semantics.

There were great works in these areas (Blok, Pigozzi, Czelakowsi, Font, Jansana, etc.), but still too detached from the specific needs of MFL.

Weakly implicative logics

Minimal reasonable behavior of an implication.

Definition

A logic L in a countable language $\mathcal{L}$ is weakly implicative if there is a binary connective $\to$ (primitive or definable) such that:

(R) $\vdash_{\mathrm{L}} \varphi \to \varphi$

(MP) $\varphi, \varphi \to \psi \vdash_{\mathrm{L}} \psi$

(T) $\varphi \to \psi, \psi \to \chi \vdash_{\mathrm{L}} \varphi \to \chi$

(sCng) $\varphi \to \psi, \psi \to \varphi \vdash_{\mathrm{L}} c(\chi_1, \dots, \chi_i, \varphi, \dots, \chi_n) \to c(\chi_1, \dots, \chi_i, \psi, \dots, \chi_n)$

for each $\langle c, n\rangle \in \mathcal{L}$ and each $0 \leq i < n$.

Such a connective is called a weak implication.

Semilinear implications and semilinear logics

A mathematical notion to capture fuzzy logics.

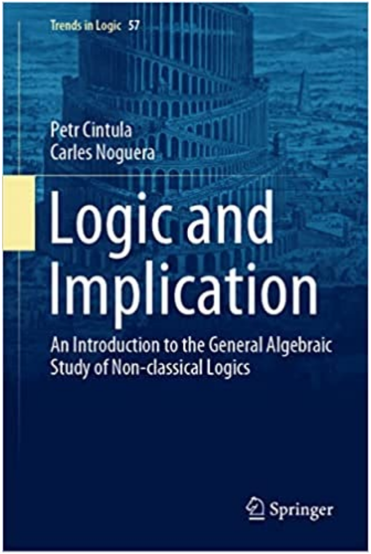
Definition

$\langle \boldsymbol{A}, F\rangle \in \mathbf{Mod}^{\ell}_{\to}(\mathrm{L})$ if for every $a, b \in A$, $a \to^{\boldsymbol{A}} b \in F$ or $b \to^{\boldsymbol{A}} a \in F$.

Definition

Let L be a logic with a weak implication $\to$. We say that L is semilinear with respect to $\to$, or alternatively that $\to$ is a semilinear implication, if $\vdash_{\mathrm{L}} = \models_{\mathbf{Mod}^{\ell}_{\to}(\mathrm{L})}$.

And now a book for the third stage



Contents

Sample of results – 1

Theorem

Let L *be a weakly implicative logic and* $\mathbb{K} \subseteq \mathbf{MOD}^*(\mathrm{L})$. *Then* L *has the* $\mathbb{K}\mathrm{C}$ *if and only if* $\mathbf{MOD}^*(\mathrm{L}) \subseteq \mathbf{HSP}(\mathbb{K})$.

Sample of results – 2

Theorem

Let L *be a weakly implicative logic and* $\mathbb{K} \subseteq \mathbf{MOD}^*(\mathrm{L})$*. Then, the following are equivalent:*

1. L *has the* $\mathrm{S}\mathbb{K}\mathrm{C}$*.*
2. $\mathbf{MOD}^*(\mathrm{L}) = \mathbf{ISP}_\omega(\mathbb{K})$*.*
3. $\mathbf{MOD}^*(\mathrm{L})^\omega \subseteq \mathbf{ISP}(\mathbb{K})$*.*

If furthermore L *has the* CIPEP*, then we can add the following equivalent condition:*

4. $\mathbf{MOD}^*(\mathrm{L})^\omega_{\mathrm{RSI}} \subseteq \mathbf{IS}(\mathbb{K})$*.*

The implication $1 \to 4$ *is always true.*

Sample of results – 3

We denote by $\mathbb{K}^+$ the extension of $\mathbb{K}$ by the trivial matrix.

Theorem

Let L *be a weakly implicative logic and* $\mathbb{K} \subseteq \mathbf{MOD}^*(\mathrm{L})$. *TFAE:*

1. L *has the* $\mathrm{FS}\mathbb{K}\mathrm{C}$.
2. $\mathbf{MOD}^*(\mathrm{L}) \subseteq \mathbf{ISPP}_{\mathrm{U}}(\mathbb{K})$.

Furthermore, if L *is finitary, then we can add:*

3. $\mathbf{MOD}^*(\mathrm{L}) = \mathbf{ISPP}_{\mathrm{U}}(\mathbb{K})$.
4. $\mathbf{MOD}^*(\mathrm{L})_{\mathrm{RFSI}}$ *is embeddable into* $\mathbf{P}_{\mathrm{U}}(\mathbb{K}^+)$, *i.e.* $\mathbf{MOD}^*(\mathrm{L})_{\mathrm{RFSI}} \subseteq \mathbf{ISP}_{\mathrm{U}}(\mathbb{K}^+)$.
5. $\mathbf{MOD}^*(\mathrm{L})^{\omega}_{\mathrm{RSI}}$ *is embeddable into* $\mathbf{P}_{\mathrm{U}}(\mathbb{K})$, *i.e.* $\mathbf{MOD}^*(\mathrm{L})^{\omega}_{\mathrm{RSI}} \subseteq \mathbf{ISP}_{\mathrm{U}}(\mathbb{K})$.

If the language of L *is finite, we can add:*

6. $\mathbf{MOD}^*(\mathrm{L})_{\mathrm{RFSI}}$ *is partially embeddable into* $\mathbb{K}^+$.
7. $\mathbf{MOD}^*(\mathrm{L})^{\omega}_{\mathrm{RSI}}$ *is partially embeddable into* $\mathbb{K}$.

Sample of results – 4

Theorem

Let L *be a weakly implicative logic with the* $\nabla\mathrm{PEP}$*. Then,*

$$\models_{\{\mathbf{B}\in\mathbf{MOD}^*(\mathrm{L})\,|\,\mathbf{B} \textit{ is linearly ordered}\}} = \mathrm{L} + (\varphi \to \psi)\ \nabla\ (\psi \to \varphi).$$

Theorem

Let L *be a logic with the* $\nabla\mathrm{PEP}$ *and let* L_1 *and* L_2 *be axiomatic extensions of* L *by sets of axioms* $\mathcal{A}_1$ *and* $\mathcal{A}_2$*, respectively. Then,* $\mathrm{L}_1 \cap \mathrm{L}_2$ *is an axiomatic extension of* L *and*

$$\mathrm{L}_1 \cap \mathrm{L}_2 = \mathrm{L} + \bigcup\{\varphi\ \nabla\ \psi \mid \varphi \in \mathcal{A}_1, \psi \in \mathcal{A}_2\}.$$

Therefore, the axiomatic extensions of L *form a sublattice of its extensions.*

Linear extension property and semilinearity property

Definition

We say that a weakly implicative logic L has the

- linear extension property, LEP for short, if linear theories form a basis of $\mathrm{Th}(\mathrm{L})$, i.e. for every theory $T \in \mathrm{Th}(\mathrm{L})$ and every formula $\varphi \in Fm_{\mathcal{L}} \setminus T$, there is a linear theory $T' \supseteq T$ such that $\varphi \notin T'$.
- semilinearity property, SLP for short, if, for each set of formulas $\Gamma \cup \{\varphi, \psi, \chi\}$,

$$\frac{\Gamma, \varphi \to \psi \vdash_{\mathrm{L}} \chi \qquad \Gamma, \psi \to \varphi \vdash_{\mathrm{L}} \chi}{\Gamma \vdash_{\mathrm{L}} \chi}.$$

Sample of results – 5

Theorem

Let L *be a weakly implicative logic. Then, the following are equivalent:*

1. L *is semilinear.*
2. L *has the* LEP*.*
3. L *has the* IPEP *and the* SLP*.*
4. L *is* RFSI*-complete and any of the following (in this context equivalent) conditions holds:*
    1. *For each* $\mathcal{L}$*-algebra* $\boldsymbol{A}$ *and each set* $X \cup \{a, b\} \subseteq A$*,* $\mathrm{Fi}(X, a \to^{\boldsymbol{A}} b) \cap \mathrm{Fi}(X, b \to^{\boldsymbol{A}} a) = \mathrm{Fi}(X)$.
    2. *Linear filters coincide with intersection-prime filters in each* $\mathcal{L}$*-algebra.*
    3. $\mathbf{MOD}^*(\mathrm{L})_{\mathrm{RFSI}} = \mathbf{MOD}^{\ell}_{\to}(\mathrm{L})$*.*

Furthermore, if L *is* RSI*-complete, we can add:*

5. $\mathbf{MOD}^*(\mathrm{L})_{\mathrm{RSI}} \subseteq \mathbf{MOD}^{\ell}_{\to}(\mathrm{L})$*.*